

Multiple-Output Functions

$x_3 \backslash x_4$	$x_1 x_2$	00	01	11	10
00				1	1
01				1	1
11		1	1		
10		1	1		

$$f_1 = x_1 x_3' + x_1' x_3 + x_2 x_3' x_4$$

$x_3 \backslash x_4$	$x_1 x_2$	00	01	11	10
00				1	1
01				1	1
11		1	1		
10		1	1		

$$f_2 = x_1 x_3' + x_1' x_3 + x_2 x_3' x_4$$

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Multiple-Output Functions

- $f = a'b'c'd + a'b'cd + a'bc'd + a'bcd + abcd + ab'c'd'$
- $g = abcd + ab'c'd' + ab'cd + ab'cd'$

cd	ab	00	01	11	10
00		0	1	1	0
01		0	1	1	0
11		0	0	1	0
10		1	0	0	0

cd	ab	00	01	11	10
00		0	0	0	0
01		0	0	0	0
11		0	0	1	0
10		1	0	1	1

- Want to implement the two functions at the same time

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First Attempt: Minimize Functions Separately

cd	ab	00	01	11	10
00		0	1	1	0
01		0	1	1	0
11		0	0	1	0
10		1	0	0	0

cd	ab	00	01	11	10
00		0	0	0	0
01		0	0	0	0
11		0	0	1	0
10		1	0	1	1

- $f = a'd + bcd + ab'c'd'$ Cost = 16 (3 AND, 1 OR, 12 inputs)
- $g = acd + ab'd'$ Cost = 11 (2 AND, 1 OR, 8 inputs)

Nothing can be shared!!!

Total cost = 27

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Better Approach: See what can be shared

cd	ab	00	01	11	10
00		0	1	1	0
01		0	1	1	0
11		0	0	1	0
10		1	0	0	0

cd	ab	00	01	11	10
00		0	0	0	0
01		0	0	0	0
11		0	0	1	0
10		1	0	1	1

- $f = a'd + abcd + ab'c'd'$
- $g = abcd + ab'c'd' + ab'c$ Only 4 distinct min terms

- Total cost = 25 (4 AND, 2 OR, 19 inputs)

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Multilevel Synthesis

- SOP or POS circuits consist of two levels (stages) of gates
- It may be possible to reduce fan-in or cost if more than 2 levels are used
- But delay may be increased

- Consider $h = abe + ce + de + abfg + cfg + dfg$
- It is already in the simplest SOP form
- Cost = 30 (6 AND, 1 OR, 23 inputs)

- By factoring,

$$h = (ab + c + d).e + (ab + c + d).fg$$

$$= (ab + c + d).(e + fg)$$
- Cost = 16 (3 AND, 2 OR, 11 inputs)

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Another Example of Multilevel Synthesis

- Consider $f = abd + cd + a'c'e + b'c'e$
- Cost = 20 (4 AND, 1 OR, 15 inputs)

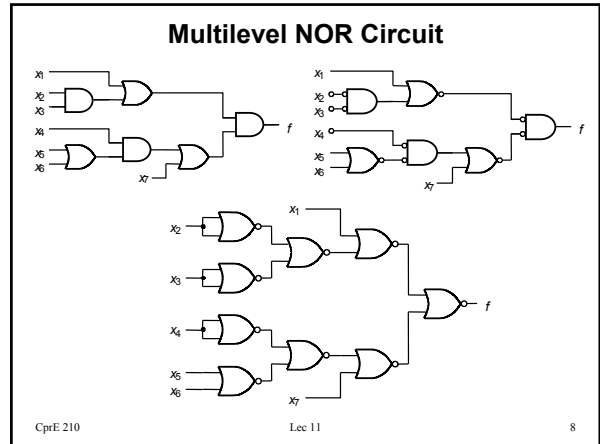
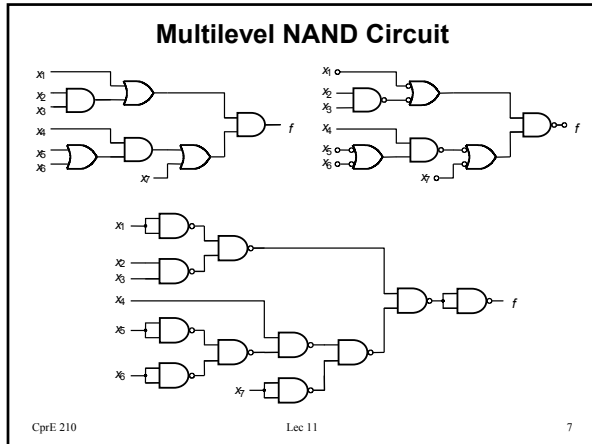
- By Boolean algebra,

$$f = (ab + c).d + (a'+b')c'e$$

$$= (ab + c).d + (ab)'c'e$$

$$= (ab + c).d + (ab + c)'e$$
- Note that the sub-function $ab + c$ appears twice in f
- We can make the implementation of f simpler by building a sub-circuit for $ab + c$.
- Cost = 17 (3 AND, 2 OR, 1 NOT, 11 inputs)

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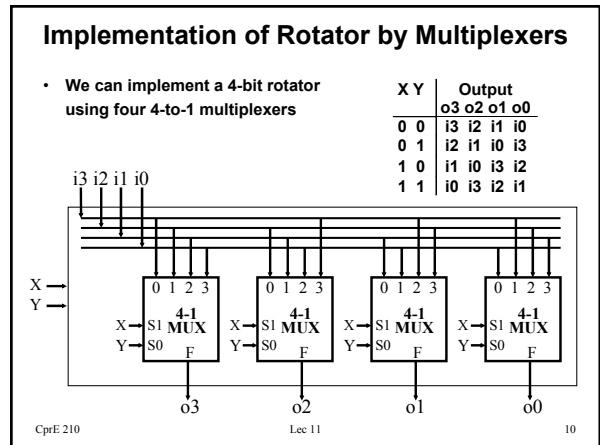


Application of Multiplexer: 4-bit rotator circuit

- Inputs are $i_3 i_2 i_1 i_0$; Outputs are $o_3 o_2 o_1 o_0$
- Output is the input rotated by 0, 1, 2 or 3 bit positions
- Rotation is to the left
- $X Y$ determines # of positions rotated to the left:
 - If $X Y = 0 0$ (left rotate 0 bit position), $o_3 o_2 o_1 o_0 = i_3 i_2 i_1 i_0$
 - If $X Y = 0 1$ (left rotate 1 bit position), $o_3 o_2 o_1 o_0 = i_2 i_1 i_0 i_3$
 - If $X Y = 1 0$ (left rotate 2 bit positions), $o_3 o_2 o_1 o_0 = i_1 i_0 i_3 i_2$
 - If $X Y = 1 1$ (left rotate 3 bit positions), $o_3 o_2 o_1 o_0 = i_0 i_3 i_2 i_1$

X Y	Output
	$o_3 o_2 o_1 o_0$
0 0	$i_3 i_2 i_1 i_0$
0 1	$i_2 i_1 i_0 i_3$
1 0	$i_1 i_0 i_3 i_2$
1 1	$i_0 i_3 i_2 i_1$

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Encoder

- Encoders perform the inverse function of Decoders.
- An encoder has 2^n (or less) input bits and n output bits
- The output bits generate the binary code corresponding to the input value
- Assuming only one input has a value of 1 at any given time
- Example: An 8-to-3 Encoder

Inputs								Outputs		
D7	D6	D5	D4	D3	D2	D1	D0	A2	A1	A0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	1	0	0	0	0	0	1	0	0
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

$A_2 = D_4 + D_5 + D_6 + D_7$
 $A_1 = D_2 + D_3 + D_6 + D_7$
 $A_0 = D_1 + D_3 + D_5 + D_7$

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Priority Encoders

- Each input signal has a priority level associated with it
- May have more than one 1's in the input signals
- Outputs indicate the active input that has the highest priority
- Example: 4-to-2 priority encoder
 - Assume w_3 has the highest priority and w_0 the lowest
 - $y_1 y_0$ indicate the active input with highest priority
 - z indicates none of the inputs is equal to 1

w3	w2	w1	w0	y1	y0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$\text{Let } i_0 = w_0 w_1' w_2' w_3'$
 $i_1 = w_1 w_2' w_3'$
 $i_2 = w_2 w_3'$
 $i_3 = w_3$
 Then $y_0 = i_1 + i_3$
 $y_1 = i_2 + i_3$
 x : both 0 and 1 (irrelevant)

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